# Time Complexity 101 and Imports

**Today starts as a Paper + Pencil or Tablet + Pencil** day... please keep laptops stowed away!

#### COMP110 - CL14 2024/03/21



### Warm-up: Diagram this Code Listing This will be today's attendance submission (CL14). Submit once complete.

1	<pre>def make_dict(keys: list[int], values: list[str]) -&gt;</pre>			
2	assert len(keys) == len(values)			
3	<pre>result: dict[int, str] = {}</pre>			
4	i: int = 0			
5	<pre>while i &lt; len(keys):</pre>			
6	result[keys[i]] = values[i]			
7	i += 1			
8	return result			
9				
10				
11	jerseys: list[int] = [2, 5, 55]			
12	<pre>names: list[str] = ["Cadeau", "Bacot", "Ingram"]</pre>			
13	<pre>players: dict[int, str] = make_dict(jerseys, names)</pre>			
14				
15	<pre>print(players[5])</pre>			
16	<pre>print(players[0])</pre>			

**Discussion question:** 

 What key difference between dictionaries and lists is illustrated in the `players` variable's value?



```
def make_dict(keys: list[int], values: list[str]) -> dict[int, str]:
         assert len(keys) == len(values)
 2
          result: dict[int, str] = {}
 3
         i: int = 0
 4
         while i < len(keys):</pre>
 5
              result[keys[i]] = values[i]
 6
              i += 1
 7
          return result
 8
 9
10
      jerseys: list[int] = [2, 5, 55]
11
     names: list[str] = ["Cadeau", "Bacot", "Ingram"]
12
      players: dict[int, str] = make_dict(jerseys, names)
13
14
      print(players[5])
15
      print(players[0])
16
```

## Warm-up #2

```
def make_dict(keys: list[int], values: list[str]) -> dict[int, str]:
          assert len(keys) == len(values)
2
          result: dict[int, str] = {}
3
          i: int = \emptyset
4
          while i < len(keys):</pre>
5
              result[keys[i]] = values[i]
6
              i += 1
          return result
8
9
10
     jerseys: list[int] = [2, 5, 55]
11
     names: list[str] = ["Cadeau", "Bacot", "Ingram"]
12
     players: dict[int, str] = make_dict(jerseys, names)
13
14
     print(players[5])
15
     print(players[0])
16
```

Suppose keys and values each have 3 elements. Assume our unit of "operation" is the number of times the block of lines #6-7 are evaluated.

**Q1.** Can different item values of keys and values lead to a difference in the number of operations required for the *intersection* function evaluation to complete?

Q2. How many operations does this function take to complete in terms of N where N is len(keys)?

### Warm-up #3: Given this example from Tuesday

1	<pre>def intersection(a: list[str], b: list[str]) -&gt; list[str</pre>
2	<pre>result: list[str] = []</pre>
3	
4	<pre>idx_a: int = 0</pre>
5	<pre>while idx_a &lt; len(a):</pre>
6	<pre>idx_b: int = 0</pre>
7	found: bool = False
8	<pre>while not found and idx_b &lt; len(b):</pre>
9	<pre>if a[idx_a] == b[idx_b]:</pre>
10	found = True
11	result.append(a[idx_a])
12	idx_b += 1
13	idx_a += 1
14	
15	return result

Suppose a and b each have 3 elements. Assume our unit of "operation" is the number of times the block of lines #9-12 are evaluated.

*intersection* function evaluation to complete? Then try for the *maximal* operations to complete? does this function take to complete?

- **Q1.** Can different values of a and b lead to a difference in the number of operations required for the
- **Q2.** If so, provide example item values for a and b which require the *fewest* operations to complete?
- Q3. Assuming the item values of a and b are random and unpredictable, about how many operations







#### Time Complexity of Common Operations of Fundamental Data Structures

#### Python's Built-in Capabilities Define Average and "Worst Case" Bounds for Common Data Structures https://wiki.python.org/moin/TimeComplexity

#### list

Operation	Average Case	C Amortized Worst Case
Сору	O(n)	O(n)
Append[1]	O(1)	O(1)
Pop last	O(1)	O(1)
Pop intermediate[2]	O(n)	O(n)
Insert	O(n)	O(n)
Get Item	O(1)	O(1)
Set Item	O(1)	O(1)
Delete Item	O(n)	O(n)
Iteration	O(n)	O(n)
Get Slice	O(k)	O(k)
Del Slice	O(n)	O(n)
Set Slice	O(k+n)	O(k+n)
Extend[1]	O(k)	O(k)
Sort Sort	O(n log n)	O(n log n)
Multiply	O(nk)	O(nk)
x in s	O(n)	
min(s), max(s)	O(n)	
Get Length	O(1)	O(1)

#### set

Operation	Average case	Worst Case
x in s	O(1)	O(n)
Union s t	O(len(s)+len(t))	
Intersection s&t	O(min(len(s), len(t)))	O(len(s) * len(t))
Multiple intersection s1&s2&&sn		(n-1)*O(I) where I is
Difference s-t	O(len(s))	
s.difference_update(t)	O(len(t))	
Symmetric Difference s^t	O(len(s))	O(len(s) * len(t))
s.symmetric_difference_update(t)	O(len(t))	O(len(t) * len(s))

#### dict

Operation	Average Case	Amortized Wo
k in d	O(1)	O(n)
Copy[3]	O(n)	O(n)
Get Item	O(1)	O(n)
Set Item[1]	O(1)	O(n)
Delete Item	O(1)	O(n)
Iteration[3]	O(n)	O(n)











Suppose a and b each had 1,000,000 elements, the worst case difference here is approximately 1,000,000 operations versus 1,000,000\*\*2 or 1,000,000,000,000 operations.

If your device can perform 100,000,000 operations per second, then...

A call to a will complete in 2.78 hours and b will complete in 1/100th of a second.

#### Orders of magnitude better...





### Imports (Code-along)

### Functions as Parameters (Code Along)

### Homework

- EXo5 Dictionary Utils Due Monday 3/25 at 11:59pm
- RDoo Ethical Algorithms Due Friday 3/29 at 11:59pm
- 3/25 at 11:59pm y 3/29 at 11:59pm